Cube Summing, Approximate Inference with Non-Local Features, and Dynamic Programming without Semirings

Kevin Gimpel and Noah A. Smith



Overview

- We introduce cube summing, which extends dynamic programming algorithms for summing with non-local features
 - □ Inspired by **cube pruning** (Chiang, 2007; Huang & Chiang, 2007)
- We relate cube summing to semiring-weighted logic programming
 - □ Without non-local features, cube summing is a novel semiring
 - □ Non-local features break some of the semiring properties
 - □ We propose an implementation based on **arithmetic circuits**



Background

- Cube Pruning
- Cube Summing
- Semirings
- Implementation
- Conclusion



Consider an exponential probabilistic model

$$p(y \mid x) \propto \prod_{m=1}^{M} \lambda_m^{h_m(x,y)}$$

Decoding

$$\hat{y}(x) = \underset{y \in \mathcal{Y}}{\operatorname{argmax}} \prod_{m=1}^{M} \lambda_m^{h_m(x,y)}$$
Summing

$$s(x) = \sum_{y \in \mathcal{Y}} \prod_{m=1}^{M} \lambda_m^{h_m(x,y)}$$



Consider an exponential probabilistic model

$$p(y \mid x) \propto \prod_{m=1}^{M} \lambda_m^{h_m(x,y)}$$

example: HMM

 ${\mathcal X}$ is a sentence, ${\mathcal Y}$ is a tag sequence

Decoding
$$\hat{y}(x) = \underset{y \in \mathcal{Y}}{\operatorname{argmax}} \prod_{m=1}^{M} \lambda_m^{h_m(x,y)}$$
Viterbi algorithm
Summing
$$s(x) = \sum_{y \in \mathcal{Y}} \prod_{m=1}^{M} \lambda_m^{h_m(x,y)}$$
forward and backward algorithms



Consider an exponential probabilistic model

$$p(y \mid x) \propto \prod_{m=1}^{M} \lambda_m^{h_m(x,y)}$$

example: PCFG x is a sentence, y is a parse tree

Decoding
$$\hat{y}(x) = \underset{y \in \mathcal{Y}}{\operatorname{argmax}} \prod_{m=1}^{M} \lambda_m^{h_m(x,y)}$$
probabilistic CKY
Summing
$$s(x) = \sum_{y \in \mathcal{Y}} \prod_{m=1}^{M} \lambda_m^{h_m(x,y)}$$
inside algorithm



Consider an exponential probabilistic model

$$p(y \mid x) \propto \prod_{m=1}^{M} \lambda_m^{h_m(x,y)}$$



Dynamic Programming

Consider the probabilistic CKY algorithm

$$egin{aligned} C_{X,i-1,i} &= \lambda_{X o w_i} \ C_{X,i,k} &= \max_{Y,Z \in \mathcal{N}; j \in \{i+1,\dots,k-1\}} \lambda_{X o YZ} imes C_{Y,i,j} imes C_{Z,j,k} \ goal &= C_{S,0,n} \end{aligned}$$



Weighted Logic Programs	Probabilistic CKY	Example
theorem	chart item	$C_{X,i,j}$
axiom	rule probability	$\lambda_{X o YZ}$
proof	derivation	PP NP of the list



Weighted Logic Programs	Probabilistic CKY	Example
theorem	chart item	$C_{X,i,j}$
axiom	rule probability	$\lambda_{X o YZ}$
proof	derivation	PP NP of the list

In semiring-weighted logic programming, theorem and axiom values come from a semiring



Features

- Recall our model: $p(y \mid x) \propto \prod_{m=1}^{M} \lambda_m^{h_m(x,y)}$
- The $h_m(x,y)$ are feature functions and the λ_m are nonnegative weights



Features

- Recall our model: $p(y \mid x) \propto \prod_{m=1}^{M} \lambda_m^{h_m(x,y)}$
- The $h_m(x,y)$ are feature functions and the λ_m are nonnegative weights
- Local features depend only on theorems used in an equation (or any of the axioms), not on the proofs of those theorems

$$C_{X,i,k} = \max_{Y,Z \in \mathcal{N}; j \in \{i+1,\dots,k-1\}} \lambda_{X \to YZ} \times C_{Y,i,j} \times C_{Z,j,k}$$











Features

- Recall our model: $p(y \mid x) \propto \prod_{m=1}^{M} \lambda_m^{h_m(x,y)}$
- The $h_m(x,y)$ are feature functions and the λ_m are nonnegative weights
- Local features depend only on theorems used in an equation (or any of the axioms), not on the proofs of those theorems

$$C_{X,i,k} = \max_{Y,Z \in \mathcal{N}; j \in \{i+1,\dots,k-1\}} \lambda_{X \to YZ} \times C_{Y,i,j} \times C_{Z,j,k}$$

Non-local features depend on theorem proofs











Other Algorithms for Approximate Inference

- Beam search (Lowerre, 1979)
- Reranking (Collins, 2000)
- Algorithms for graphical models
 - □ Variational methods (MacKay, 1997; Beal, 2003; Kurihara & Sato, 2006)
 - □ Belief propagation (Sutton & McCallum, 2004; Smith & Eisner, 2008)
 - □ MCMC (Finkel et al., 2005; Johnson et al., 2007)
 - □ Particle filtering (Levy et al., 2009)
- Integer linear programming (Roth & Yih, 2004)
- Stacked learning (Cohen & Carvalho, 2005; Martins et al., 2008)
- Cube pruning (Chiang, 2007; Huang & Chiang, 2007)



Other Algorithms for Approximate Inference

- Beam search (Lowerre, 1979)
- Reranking (Collins, 2000)
- Algorithms for graphical models
 - □ Variational methods (MacKay, 1997; Beal, 2003; Kurihara & Sato, 2006)
 - □ Belief propagation (Sutton & McCallum, 2004; Smith & Eisner, 2008)
 - □ MCMC (Finkel et al., 2005; Johnson et al., 2007)
 - □ Particle filtering (Levy et al., 2009)
- Integer linear programming (Roth & Yih, 2004)
- Stacked learning (Cohen & Carvalho, 2005; Martins et al., 2008)
- Cube pruning (Chiang, 2007; Huang & Chiang, 2007)
- Why add one more?
 - Cube pruning extends existing, widely-understood dynamic programming algorithms for decoding
 - □ We want this for **summing** too



- Background
- Cube Pruning
- Cube Summing
- Semirings
- Implementation
- Conclusion



Cube Pruning (Chiang, 2007; Huang & Chiang, 2007)

- Modification to dynamic programming algorithms for decoding to use non-local features approximately
- Keeps a k-best list of proofs for each theorem
- Applies non-local feature functions on these proofs when proving new theorems





$$C_{NP,0,7} = C_{NP,0,1} \times C_{PP,1,7} \times \lambda_{NP \to NP PP}$$





$$C_{NP,0,7} = C_{NP,0,1} \times C_{PP,1,7} \times \lambda_{NP \rightarrow NP PP}$$





$$C_{NP,0,7} = C_{NP,0,1} \times C_{PP,1,7} \times \lambda_{NP \to NP PP}$$





$$C_{_{NP,0,7}} = C_{_{NP,0,1}} \times C_{_{PP,1,7}} \times \lambda_{_{NP \to NP PP}}$$

























Clarification

- Cube pruning does not actually expand all k² proofs as this example showed
- It uses an approximation that only looks at O(k) proofs
- But since we are summing, we want to look at as many proofs as possible
- We use the algorithm that we just showed as the basis for cube summing (we call it cube decoding – details in paper)



- Background
- Cube Pruning
- Cube Summing
- Semirings
- Implementation
- Conclusion











- Computation of local and non-local features is same as before
- Only difference is computing the residual for the result

	C				
С	- PP,1,7 NP,0,1	0.2	0.1	0.05	0.03
	0.4	0.008	0.004	0.001	
	0.3	0.018	0.009	0.003	
	0.02	0.0002	0.0001	0.0001	
	0.05				





	C_{DD}				
C	PP,1,7 NP,0,1	0.2	0.1	0.05	0.03
	0.4	0.008			
	0.3	0.018	0.009	0.0004	
	0.02			0.0084	
	0.05				











 $C_{NP.0.7} = C_{NP.0.1} \times C_{PP.1.7} \times \lambda_{NP \to NP PP}$

 $\lambda_{_{NP \rightarrow NP PP}} = 0.5$



C _{NP,0,7}	0.018	0.009	0.008	0.0287	
----------------------------	-------	-------	-------	--------	--



C	- PP,1,7 NP,0,1	0.2	0.1	0.05	0.03
	0.4	0.008			
	0.3	0.018	0.009	0.0004	
	0.02			0.0084	
	0.05	0.005	0.0025	0.00125	





	C				
C	⁻ PP,1,7 NP,0,1	0.2	0.1	0.05	0.03
	0.4	0.008			
	0.3	0.018	0.009	0.0004	
	0.02			0.0084	
	0.05		0.00875		





	C_{DDAZ}				
C	PP,1,7 NP,0,1	0.2	0.1	0.05	0.03
	0.4	0.008			0.012 × 0.5
	0.3	0.018	0.009	0.0084	0.009 × 0.5
	0.02				0.0006 × 0.5
	0.05		0.00875		
•					





	C_{PP17}				
C_{i}	PP,1,7 NP,0,1	0.2	0.1	0.05	0.03
	0.4	0.008			
	0.3	0.018	0.009	0.0004	0.0108
	0.02			0.0084	
	0.05		0.00875		





	C_{DD}				
C	- PP,1,7 NP,0,1	0.2	0.1	0.05	0.03
	0.4	0.008			
	0.3	0.018	0.009	0.0004	0.0108
	0.02			0.0084	
	0.05		0.00875		0.0015 × 0.5





	$C_{\text{DD4.7}}$				
C	PP,1,7 NP,0,1	0.2	0.1	0.05	0.03
	0.4	0.008			
	0.3	0.018	0.009	0.0004	0.0108
	0.02			0.0084	
	0.05		0.00875		0.00075







С _{NP,0,7}	0.018	0.009	0.008	0.0287	
---------------------	-------	-------	-------	--------	--



Summary

- Maintain residual sum of all proofs not in k-best list
- Redefine operations to update the residual as necessary
- Result is approximate k-best proof list for goal and approximate sum of all other proofs of goal
- When $k = \infty$, result is exact



- Background
- Cube Pruning
- Cube Summing
- Semirings
- Implementation
- Conclusion



Semirings

Semiring	A	\oplus	\otimes	0	1
Inside	$\mathbb{R}_{\geq 0}$	a+b	ab	0	1
Viterbi	$\mathbb{R}_{\geq 0}$	$\max(a,b)$	ab	0	1

• A semiring is a tuple $\langle A, \oplus, \otimes, \mathbf{0}, \mathbf{1} \rangle$ such that: $\square \oplus : A \times A \to A$ is associative and commutative $\square \otimes : A \times A \to A$ is associative and distributes over \bigoplus $\square \forall a \in A, a \oplus \mathbf{0} = a,$ $a \otimes \mathbf{1} = a,$ $a \otimes \mathbf{0} = \mathbf{0} \otimes a = \mathbf{0}$



Non-local features break some of the semiring properties! (see paper for details)





- Background
- Cube Pruning
- Cube Summing
- Semirings
- Implementation
- Conclusion



Implementation

- Several implementation tools exist for dynamic programming
 - Dyna (Eisner et al., 2005) and Goodman (1999) assume semirings
 - Hypergraphs (Klein & Manning, 2001; Huang, 2008) do not require semirings but are aimed at decoding
- These could be extended for cube summing, but we instead use a lower-level formalism: arithmetic circuits



Arithmetic Circuits

- Explicitly represent computations to be performed using a directed graph
 - Operators and operands are nodes in the graph
 - □ A value is associated with each node
 - Operators point to their operands
- Allow automatic differentiation in the reverse mode (Griewank & Corliss, 1991) for efficient gradient computation





- Background
- Cube Pruning
- Cube Summing
- Semirings
- Implementation
- Conclusion



Conclusion and Ongoing Work

- We have described cube summing, a technique for approximate summing using dynamic programming with non-local features
- With only local features, cube summing is a semiring that generalizes those in common use
- Some semiring properties are broken by non-local features but an implementation based on arithmetic circuits can be used
- We are currently using cube summing to train a loglinear syntactic translation model with hidden variables



Thanks!

Cube Summing, Approximate Inference with Non-Local Features, and Dynamic Programming without Semirings

Kevin Gimpel and Noah A. Smith

