# An $\mathcal{O}(n \log n)$ projection operator for weighted $\ell_{1}$-norm regularization with sum constraint 

Weiran Wang<br>Toyota Technological Institute at Chicago<br>weiranwang@ttic.edu

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#### Abstract

We provide a simple and efficient algorithm for the projection operator for weighted $\ell_{1}$-norm regularization subject to a sum constraint, together with an elementary proof. The implementation of the proposed algorithm can be downloaded from the author's homepage.


## 1 The problem

In this report, we consider the following optimization problem:

$$
\begin{array}{ll}
\min _{\mathbf{x}} & \frac{1}{2}\|\mathbf{x}-\mathbf{y}\|^{2}+\sum_{i=1}^{n} d_{i}\left|x_{i}\right|  \tag{1}\\
\text { s.t. } & \mathbf{x}^{\top} \mathbf{1}=1
\end{array}
$$

where $\mathbf{y}=\left[y_{1}, \ldots, y_{n}\right]^{\top} \in \mathbb{R}^{n}, d_{i} \geq 0, i=1, \ldots, n$, and $\mathbf{1}$ is the $n$-dimensional vector consisting of all 1 's. This is a quadratic program and the objective function is strictly convex (even though it is non-smooth), so there is a unique solution which we denote by $\mathbf{x}=\left[x_{1}, \ldots, x_{n}\right]^{\top}$ with a slight abuse of notation.

Notice if $d_{1}=\cdots=d_{n}$ and the constraint were absent, the problem has a closed form solution known as the soft-shrinkage operator (see, e.g., Beck and Teboulle, 2009), which is widely used for solving $\ell_{1}$-regularized problem in learning sparse representations. But our problem (1) is more involved due to the constraint that couples all dimensions of $\mathbf{x}$. Nonetheless, we give an efficient algorithm with time complexity $\mathcal{O}(n \log n)$ for this problem using only the KKT theorem.
Remark 1.1. Our motivation for (1) also comes from sparse coding. Yu et al. (2009) propose the local coordinate coding (LCC) algorithm for learning sparse representations induced by locality. Given a data sample $\mathbf{u} \in \mathbb{R}^{n}$ and a set of landmark points $\left\{\mathbf{v}_{j}\right\}_{j=1}^{C}$ where $\mathbf{v}_{j} \in \mathbb{R}^{n}, j=1, \ldots, C$, the LCC algorithm reconstructs $\mathbf{u}$ from the landmark points while enforcing the faraway landmark points to contribute less than nearby landmark points (or to have smaller reconstruction coefficients). Let the reconstruction coefficient of $\mathbf{v}_{j}$ be $w_{j}, j=1, \ldots, C$. Then the optimization problem for these coefficients in LCC is

$$
\begin{array}{ll}
\underset{\mathbf{w}}{\min } & \left\|\mathbf{u}-\sum_{j=1}^{C} w_{j} \mathbf{v}_{j}\right\|^{2}+\lambda \sum_{j=1}^{C}\left\|\mathbf{u}-\mathbf{v}_{j}\right\|^{2}\left|w_{j}\right|  \tag{2}\\
\text { s.t. } & \sum_{j=1}^{C} w_{j}=1
\end{array}
$$

where $\lambda>0$ is some trade-off parameter. The constraint in (2) ensures that the representation is translation invariant. There are different ways of solving this problem, e.g., Elhamifar and Vidal (2011) have a similar optimization problem which they solve with Alternating Direction Method of Multipliers (Boyd et al., 2011).

One simple way of solving (2) is to use the gradient proximal algorithm and its Nesterov's acceleration scheme (see Beck and Teboulle, 2009 and the reference therein), where one iteratively takes a short gradient step for the smooth quadratic term and projects the new estimate with the weighted $\ell_{1}$ regularization term subject to the sum constraint, where the projection operator solves exactly (1).

## 2 The solution

We solve the problem (1) using only the KKT theorem (Nocedal and Wright, 2006), which states the necessary and sufficient condition ${ }^{1}$ satisfied by the solution $\mathbf{x}$. The Lagrangian of (1) is

$$
\begin{equation*}
\mathcal{L}(\mathbf{x}, \alpha)=\frac{1}{2}\|\mathbf{x}-\mathbf{y}\|^{2}+\sum_{i=1}^{n} d_{i}\left|x_{i}\right|+\alpha\left(\mathbf{x}^{\top} \mathbf{1}-1\right) \tag{3}
\end{equation*}
$$

where $\alpha$ is the Lagrange multipliers associated with the constraint. And the KKT system of this problem is

$$
\begin{gather*}
x_{i}-y_{i}+d_{i}+\alpha=0, \quad \text { if } \quad x_{i}>0  \tag{4a}\\
x_{i}-y_{i}-d_{i}+\alpha=0, \quad \text { if } \quad x_{i}<0  \tag{4b}\\
-d_{i} \leq-y_{i}+\alpha \leq d_{i},  \tag{4c}\\
\text { if } \quad x_{i}=0  \tag{4~d}\\
\sum_{i=1}^{n} x_{i}=1,
\end{gather*}
$$

where we have used the fact that the sub-differential of $|x|$ is $[-1,1]$ at $x=0$ to obtain (4c).
Denote $y_{i}^{-}=y_{i}-d_{i}, y_{i}^{+}=y_{i}+d_{i}, i=1, \ldots, n$, which can be computed beforehand. We can then rewrite (4) in terms of $\alpha$ :

$$
\begin{align*}
& \alpha<y_{i}^{-} \Longleftrightarrow x_{i}>0  \tag{5a}\\
& \alpha>y_{i}^{+} \Longleftrightarrow x_{i}<0  \tag{5b}\\
& y_{i}^{-} \leq \alpha \leq y_{i}^{+} \Longleftrightarrow x_{i}=0  \tag{5c}\\
& \sum_{i: x_{i}>0}\left(y_{i}^{-}-\alpha\right)+\sum_{i: x_{i}<0}\left(y_{i}^{+}-\alpha\right)=1 \tag{5~d}
\end{align*}
$$

Obviously, the Lagrange multiplier $\alpha$ is the key to our problem. Once the value of $\alpha$ is determined, we can easily obtain the optimal solution by setting

$$
\begin{array}{ll}
x_{i}=y_{i}^{-}-\alpha & \text { if } y_{i}^{-}>\alpha \\
x_{i}=y_{i}^{+}-\alpha & \text { if } y_{i}^{+}<\alpha \\
x_{i}=0 & \text { otherwise } \tag{6c}
\end{array}
$$

We can sort all dimensions of $y_{i}^{-}$and $y_{i}^{+}$together (a total of $2 N$ scalars) into an ascending $z$-sequence:

$$
\begin{equation*}
z_{1} \leq z_{2} \leq \cdots \leq z_{2 N} \tag{7}
\end{equation*}
$$

An important observation is that the $z$-sequence partitions the real axis into $4 N+1$ disjoint sets, each being either a single point set $\left\{z_{j}\right\}, j=1, \ldots, 2 N$ or an open interval of the form $\left(-\infty, z_{1}\right),\left(z_{j}, z_{j+1}\right)$, $j=1, \ldots, 2 N-1$, or $\left(z_{2 N}, \infty\right)$ and the Lagrange multiplier $\alpha$ for the solution must lie in one of them.

We then test each of the $4 N+1$ sets as follows. Assuming that $\alpha$ lies in one set, we can use (5a)-(5c) to conjecture the positive, negative, and zero dimensions of a possible solution $\hat{\mathbf{x}}$. After that, we use (5d) to compute a hypothesized value $\hat{\alpha}$ for the Lagrange multiplier, i.e.,

$$
\begin{equation*}
\alpha=\frac{\sum_{i: \hat{x}_{i}>0} y_{i}^{-}+\sum_{i: \hat{x}_{i}<0} y_{i}^{+}-1}{\sum_{i: \hat{x}_{i}>0} 1+\sum_{i: \hat{x}_{i}<0} 1} . \tag{8}
\end{equation*}
$$

[^0]If the computed $\hat{\alpha}$ indeed lies in the assumed set (a point or an open interval), we have a KKT point and thus the solution.

Since the problem (1) is strictly convex and there exists a unique global optimum, this procedure will find the exact solution with no more than $4 N+1$ tests. We can do this efficiently by sorting $y_{i}^{-}$and $y_{i}^{+}$separately $(\mathcal{O}(n \log n)$ operations) and gradually merging the two sorted sequences (an $\mathcal{O}(n)$ operation). Therefore the total cost of our procedure for solving (1) is of order $\mathcal{O}(n \log n)$.

Algorithm 1 gives the detailed pseudocode for solving (1), whose MATLAB and C++ implementation can be downloaded at https://eng.ucmerced.edu/people/wwang5.

## References

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Algorithm 1 Pseudo-code of our projection operator for (1).
Input: \(\mathbf{y} \in \mathbb{R}^{n}\) and \(\mathbf{d}=\left[d_{1}, \ldots, d_{n}\right]\) where \(d_{i} \geq 0, i=1, \ldots, n\).
    Sort \(\mathbf{y}-\mathbf{d}\) into \(\mathbf{y}^{-}: y_{1}^{-} \leq y_{2}^{-} \leq \cdots \leq y_{n}^{-}\). And sort \(\mathbf{y}+\mathbf{d}\) into \(\mathbf{y}^{+}: y_{1}^{+} \leq y_{2}^{+} \leq \cdots \leq y_{n}^{+}\).
    \(i \leftarrow 1, j \leftarrow 1 \quad \% i / j\) index of the dimension of \(\mathbf{y}^{-} / \mathbf{y}^{+}\)that will be merged next.
    \(\% s_{1} / s_{2}\) stores the sum of dimensions of \(\mathbf{y}^{-} / \mathbf{y}^{+}\)that are strictly greater/smaller than the current estimate of \(\alpha\).
    \(s_{1} \leftarrow \sum_{i=1}^{n} y_{i}^{-}, s_{2} \leftarrow 0, t \leftarrow n \quad \% t\) is the number of nonzero dimensions of the hypothesized \(\mathbf{x}\).
    if \((s 1+s 2)<t \cdot y_{1}^{-}\)then
        \(\alpha \leftarrow(s 1+s 2) / t\); return \(\quad \% \alpha<y_{1}^{-}\), all dimensions of \(\mathbf{x}\) are positive.
    end if
    while true do
        \% Test a single point set.
        if \(y_{i}^{-}<y_{j}^{+}\)then
            \(k \leftarrow i \quad \% y_{i}^{-}\)is the next value in the \(z\)-sequence.
            while \(\left(y_{k}^{-}=y_{i}^{-}\right) \& \&(k \leq n)\) do
                \(s_{1} \leftarrow s_{1}-y_{k}^{-}, t \leftarrow t-1, k \leftarrow k+1 \quad\) \% Skip the contiguous block of identical dimensions in \(\mathbf{y}^{-}\).
            end while
            if \(\left(s_{1}+s_{2}-1\right)=t \cdot y_{i}^{-}\)then
                \(\alpha \leftarrow y_{i}^{-} ;\)return \(\quad \% \alpha\) happens to lie in a single point set.
            else
                left \(\leftarrow y_{i}^{-}, i \leftarrow k \quad\) \% Otherwise, \(\alpha\) lies in a open interval with left boundary left.
            end if
        else
            if \(y_{i}^{-}>y_{j}^{+}\)then
                \(\% y_{j}^{+}\)is the next value in the \(z\)-sequence.
                    if \(\left(s_{1}+s_{2}-1\right)=t \cdot y_{j}^{+}\)then
                \(\alpha \leftarrow y_{j}^{+}\); return \(\quad \% \alpha\) happens to lie in a single point set.
                    else
                                left \(\leftarrow y_{j}^{+} \quad \%\) Otherwise, \(\alpha\) lies in a open interval with left boundary left.
                while \(\left(y_{j}^{+}=l e f t\right) \& \&(j \leq n)\) do
                        \(s_{2} \leftarrow s_{2}+y_{j}^{+}, t \leftarrow t+1, j \leftarrow j+1 \quad\) \% Skip the contiguous block of identical entries in \(\mathbf{y}^{+}\).
                end while
                    end if
            else
                \(k \leftarrow i \quad \% y_{i}^{-}=y_{j}^{+}\)is the next value in the \(z\)-sequence.
                    while \(\left(y_{k}^{-}=y_{i}^{-}\right) \& \&(k \leq n)\) do
                            \(s_{1} \leftarrow s_{1}-y_{k}^{-}, t \leftarrow t-1, k \leftarrow k+1\)
                end while
                if \(\left(s_{1}+s_{2}-1\right)=t \cdot y_{i}^{-}\)then
                    \(\alpha \leftarrow y_{i}^{-}\); return
                    else
                    \(l e f t \leftarrow y_{i}^{-}, i \leftarrow k\)
                    while \(\left(y_{j}^{+}=l e f t\right) \& \&(j \leq n)\) do
                            \(s_{2} \leftarrow s_{2}+y_{j}^{+}, t \leftarrow t+1, j \leftarrow j+1\)
                    end while
                end if
            end if
        end if
        \% Find the right boundary of the open interval and test if it contains \(\alpha\).
        if \(y_{i}^{-}<y_{j}^{+}\)then
            right \(\leftarrow y_{i}^{-}\)
        else
            right \(\leftarrow y_{j}^{+}\)
        end if
        if \(t \cdot\) left \(<\left(s_{1}+s_{2}-1\right) \& \& t \cdot\) right \(>\left(s_{1}+s_{2}-1\right)\) then
            \(\alpha \leftarrow\left(s_{1}+s_{2}-1\right) / t\); return
                        \(\% \alpha\) lies in the open interval (left, right).
        end if
    end while
Output: \(\alpha\) is the Lagrange multiplier of the problem (1), use (6) to obtain \(\mathbf{x}\).
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[^0]:    ${ }^{1}$ Strictly speaking, our objective is convex and non-smooth, so the condition is that the zero vector $\mathbf{0}$ lies in the sub-differential at the solution $\mathbf{x}$.

