Information and Coding Theory

Homework 2

Due: February 12, 2021

1. Biased coins strike back.

[3+3=6 points]

Winter 2021

In class we considered the problem of distinguishing coins distributed according to the following two distributions:

$$P = \begin{cases} 1 & \text{w.p. } \frac{1}{2} - \varepsilon \\ 0 & \text{w.p. } \frac{1}{2} + \varepsilon \end{cases} \text{ and } Q = \begin{cases} 1 & \text{w.p. } \frac{1}{2} \\ 0 & \text{w.p. } \frac{1}{2} \end{cases}$$

We derived matching upper and lower bounds (up to constants) of the form $\Theta(1/\epsilon^2)$ on the number of coin tosses required to distinguish the two distributions. Consider now the problem of distinguishing two extremely biased coins with slightly differing biases:

$$P' = \begin{cases} 1 & \text{w.p. } \varepsilon \\ 0 & \text{w.p. } 1 - \varepsilon \end{cases} \text{ and } Q' = \begin{cases} 1 & \text{w.p. } 2\varepsilon \\ 0 & \text{w.p. } 1 - 2\varepsilon \end{cases}$$

Find tight upper and lower bounds (up to constants) on the number of independent coin tosses required to distinguish coins distributed according to P' and Q'.

2. Jensen-Shannon divergence.

[2+3+4+3=12 points]

While KL-divergence is sometimes used as a measure of the difference between two distributions, it is asymmetric and can be infinite. In some applications, one can instead consider the Jensen-Shanon divergence which addresses these issues.

(a) For two distributions *P* and *Q*, we define the Jensen-Shannon divergence as

$$JSD(P,Q) := \frac{1}{2} \cdot D(P||M) + \frac{1}{2} \cdot D(Q||M) \text{ where } M = \frac{P+Q}{2}.$$

Show that $0 \leq \text{JSD}(P, Q) \leq 1$.

- (b) Show that $JSD(P,Q) \ge \frac{1}{8 \ln 2} \cdot ||P Q||_1^2$.
- (c) Show that $JSD(P,Q) \leq \frac{1}{2} \cdot ||P-Q||_1$.
- (d) The notion of Jensen-Shannon divergence can be generalized to an arbitrary number of distributions and an arbitrary convex combination. Let P_1, \ldots, P_k be distributions on the same universe and let $\lambda = (\lambda_1, \ldots, \lambda_k)$ be a tuple of non-negative weights such that $\sum_i \lambda_i = 1$. We define

$$JSD_{\lambda}(P_1,\ldots,P_k) := \sum_i \lambda_i \cdot D(P_i || M) \text{ where } M = \sum_i \lambda_i P_i.$$

Show that $0 \leq \text{JSD}_{\lambda}(P_1, \dots, P_k) \leq H(\lambda)$, where $H(\lambda)$ denotes the entropy of λ , when viewed as a distribution over [k].

3. Counting using method of types (Problem 11.5 from the book). [5 points] Let \mathcal{X} be a finite universe with $|\mathcal{X}| = r$ and let $g : \mathcal{X} \to \mathbb{R}$ be a real valued function. Let $S \subseteq \mathcal{X}^n$ be the set of sequences x_1, \ldots, x_n with each $x_i \in \mathcal{X}$ defined as

$$S = \left\{ (x_1,\ldots,x_n) \in \mathcal{X}^n \mid \frac{1}{n} \sum_{i=1}^n g(x_i) \ge \alpha \right\}.$$

Let $\Pi = \{P \mid \sum_{a \in \mathcal{X}} P(a)g(a) \ge \alpha\}$. Show that

$$|S| \leq (n+1)^r \cdot 2^{nH^*}$$
,

where $H^* = \sup_{P \in \Pi} H(P)$.

4. Differential entropy of a Gaussian.

[2 + 3 = 5 points]

We saw in class that if the differential entropy h(X) exists for a continuous random variable X taking values in \mathbb{R}^n , and $A \in \mathbb{R}^{n \times n}$ is a non-singular matrix, then

$$h(AX) = h(X) + \log|A|,$$

where |A| denotes |Det(A)|. We can use this to compute the entropy of a Gaussian random variable.

(a) Let $X \sim N(\mu, \Sigma)$ be an *n*-dimensional Gaussian with mean μ and covariance matrix Σ i.e.,

$$\mathbb{E}[X] = \mu$$
 and $\mathbb{E}\left[(X - \mu)(X - \mu)^{\mathsf{T}}\right] = \Sigma$.

Assume that the covariance matrix Σ is *positive definite* and hence there exists a non-singular matrix R such that $\Sigma = R^2$. Use this to show that

$$h(X) = \frac{n}{2} \cdot \log(2\pi e) + \frac{1}{2} \cdot \log|\Sigma| .$$

(b) Use the above to show that for any two positive definite matrices Σ_1 and Σ_2 , and $\alpha \in [0, 1]$, we have

$$|\alpha \cdot \Sigma_1 + (1-\alpha) \cdot \Sigma_2| \geq |\Sigma_1|^{\alpha} \cdot |\Sigma_2|^{1-\alpha}$$

5. Chernoff bound for read-*k* families.

[3+3+3+3 = 12 points]

We used Sanov's theorem to derive the Chernoff bound for independent random variables X_1, \ldots, X_n taking values uniformly in $\{0, 1\}$. In particular, we showed that

$$\mathbb{P}\left[X_1+\cdots+X_n\geq \left(\frac{1}{2}+\varepsilon\right)n\right] \leq (n+1)\cdot 2^{-n\cdot D\left(\frac{1}{2}+\varepsilon\|\frac{1}{2}\right)},$$

where $D\left(\frac{1}{2} + \varepsilon \| \frac{1}{2}\right)$ denotes the KL-divergence of two distributions on $\{0,1\}$, with probabilities $\left(\frac{1}{2} + \varepsilon, \frac{1}{2} - \varepsilon\right)$ and $\left(\frac{1}{2}, \frac{1}{2}\right)$. In this problem, we will consider functions f_1, \ldots, f_r depending on the variables X_1, \ldots, X_n and prove a concentration bound on the expression $f_1 + \cdots + f_r$.

Let S_1, \ldots, S_r be subsets of [n] for each $i \in [r]$, let $f_i : \{0,1\}^{S_i} \to \{0,1\}$ be a function which depends only on the variables in S_i . We use the shorthand X_{S_i} to denote the variables $\{X_j\}_{j \in S_i}$. Moreover, we have the property that each variable is involved in only k functions i.e., $\forall j \in [n]$, $|\{i \in [r] \mid j \in S_i\}| = k$. Such a family of functions is called a read-k family (it is not too hard to see that the lower bound extends to the case when each variable is in *at most* k functions).

(a) Recall that for two random variables Z_1 and Z_2 distributed on *same universe* Z with distributions P_1 and P_2 , we also use $D(Z_1||Z_2)$ to mean $D(P_1||P_2)$. Let Y_1, \ldots, Y_n be (not necessarily independent) random variables jointly distributed on $\{0,1\}^n$ and let X_1, \ldots, X_n be random variables as above, distributed uniformly and independently on $\{0,1\}^n$. Let the sets $\{S_i\}_{i \in [r]}$ be as above. Use Shearer's lemma to show that

$$k \cdot D(Y_1, \ldots, Y_n || X_1, \ldots, X_n) \geq \sum_{i \in [r]} D(Y_{S_i} || X_{S_i}).$$

(b) Let $A = \{(a_1, \ldots, a_n) \in \{0, 1\}^n \mid \sum_{i \in [r]} f_i(\{a_j\}_{j \in S_i}) \ge t\}$. Let (Y_1, \ldots, Y_n) be uniformly distributed over the set A (note that Y_1, \ldots, Y_n are not necessarily independent). Prove that

$$\mathbb{P}_{X_{1},...,X_{n}}\left[\sum_{i\in[r]}f_{i}(X_{S_{i}})\geq t\right] = 2^{-D(Y_{1},...,Y_{n}||X_{1},...,X_{n})},$$

where the probability is over the uniform distribution for X_1, \ldots, X_n .

(c) For each $i \in [r]$, let $\mathbb{E}[f_i(X_{S_i})] = \mu_i$ and $\mathbb{E}[f_i(Y_{S_i})] = \nu_i$. Prove that

$$D(Y_{S_i} \| X_{S_i}) \geq D(\nu_i \| \mu_i)$$
 ,

where $D(v_i || \mu_i)$ denotes the divergence of two distributions on $\{0, 1\}$ with probabilities $(v_i, 1 - v_i)$ and $(\mu_i, 1 - \mu_i)$.

(d) Use the above bounds and the convexity of KL-divergence in both its arguments to show that for $\mu = \frac{1}{r} \cdot (\mu_1 + \cdots + \mu_r)$,

$$\mathbb{P}_{X_1,\ldots,X_n}\left[f_1(X_{S_1})+\cdots+f_r(X_{S_r})\geq (\mu+\varepsilon)\cdot r\right] \leq 2^{-(r/k)\cdot D(\mu+\varepsilon\|\mu)}.$$